Predicate Calculus Validity

Propositional validity

$$
(A \to B) \lor (B \to A)
$$

True *no matter what* the truth values of *A* and *B* are

Predicate calculus validity

 $\forall z \left[Q(z) \wedge P(z) \right] \rightarrow \left[\forall x . Q(x) \wedge \forall y . P(y) \right]$

True *no matter what*

- the Domain is,
- or the predicates are.

That is, logically correct, independent of the specific content.

Arguments with Quantified Statements

Universal instantiation:

$$
\mathop{\qquad \forall x, P(x)}\\ \cdot \hspace{0.1cm} P(a)
$$

Universal modus ponens:

 \bullet

 \bullet

$$
\forall x, P(x) \to Q(x)
$$

$$
P(a)
$$

$$
\cdot Q(a)
$$

Universal modus tollens:

$$
\forall x, P(x) \to Q(x)
$$

$$
\neg Q(a)
$$

$$
\cdot \neg P(a)
$$

Universal Generalization

 (c) $R(x)$ $A \rightarrow R(c)$ $A \rightarrow \forall x.R(x)$ \rightarrow $\rightarrow \forall x$ valid rule

providing *c* is independent of *A*

Informally, if we could prove that R(c) is true for an arbitrary c (in a sense, c is a "variable"), then we could prove the for all statement.

e.g. given any number c, 2c is an even number

 \Rightarrow for all x, 2x is an even number.

Remark: Universal generalization is often difficult to prove, we will introduce mathematical induction to prove the validity of for all statements.

Valid Rule?

 $\forall z \left[\mathbb{Q}(z) \mathsf{V} \mathbb{P}(z)\right] \rightarrow \left[\forall x \mathbb{Q}(x) \mathsf{V} \forall y \mathbb{P}(y)\right]$

Proof: Give *countermodel*, where $\forall z$ $[Q(z)$ **V** $P(z)$] is true, but $\forall x. Q(x)$ **V** $\forall y. P(y)$ is false.

Find a domain, and a predicate.

In this example, let domain be integers, *Q*(*z*) be true if z is an even number, i.e. Q(z)=even(z) *P*(*z*) be true if z is an odd number, i.e. P(z)=odd(z)

Then $\forall z$ [Q(*z*) **V** $P(z)$] is true, because every number is either even or odd. But $\forall x. Q(x)$ is not true, since not every number is an even number. Similarly $\forall y. P(y)$ is not true, and so $\forall x. Q(x)$ **V** $\forall y. P(y)$ is not true.

Valid Rule?

$\forall z \in D$ $[Q(z) \wedge P(z)] \rightarrow [\forall x \ Q(x) \wedge \forall y \ P(y)]$

Proof: Assume $\forall z$ [*Q*(*z*) \land *P*(*z*)].

So $Q(z)/P(z)$ holds for all *z* in the domain D.

Now let *c* be some element in the domain D.

So *Q*(*c*) *P*(*c*) holds (by instantiation), and therefore *Q*(*c*) by itself holds.

But *c* could have been any element of the domain D.

So we conclude $\forall x. Q(x)$. (by generalization)

We conclude $\forall y.P(y)$ similarly (by generalization). Therefore,

 $\forall x. Q(x) \land \forall y. P(y)$ QED.

This Lecture

Now we have learnt the basics in logic.

We are going to apply the logical rules in proving mathematical theorems.

- Direct proof
- Contrapositive
- Proof by contradiction
- Proof by cases

Basic Definitions

An integer n is an even number

if there exists an integer k such that $n = 2k$.

An integer n is an odd number

if there exists an integer k such that $n = 2k+1$.

Proving an Implication

Goal: If P, then Q. (P implies Q)

Method 1: Write assume P, then show that Q logically follows.

$$
\boxed{\text{Claim:}} \qquad \text{If } 0 \le x \le 2 \text{, then } -x^3 + 4x + 1 > 0
$$

Reasoning: When $x=0$, it is true.

When x grows, 4x grows faster than x^3 in that range.

Proof:
$$
-x^3 + 4x + 1 = x(2-x)(2+x) + 1
$$

When $0 \le x \le 2$, $x(2-x)(2+x) \ge 0$

Direct Proofs

The sum of two even numbers is even.

Proof

\n
$$
x = 2m, y = 2n
$$
\n
$$
x+y = 2m+2n
$$
\n
$$
= 2(m+n)
$$

The product of two odd numbers is odd.

x = 2m+1, y = 2n+1 Proof $xy = (2m+1)(2n+1)$ $= 4mn + 2m + 2n + 1$ $= 2(2mn+mn+n) + 1.$

Divisibility

a "divides" b (a|b):

b = ak for some integer k

5|15 because 15 = 3X5 n|0 because 0 = nX0 1|n because n = 1Xn n|n because n = nX1

A number $p > 1$ with no positive integer divisors other than 1 and itself is called a **prime**. Every other number greater than 1 is called **composite**.

2, 3, 5, 7, 11, and 13 are prime,

4, 6, 8, and 9 are composite.

Simple Divisibility Facts

```
1. If a |b, then a |bc for all c.
2. If a \vert b and b \vert c, then a \vert c.
3. If a \vert b and a \vert c, then a \vert sb + tc for all s and t.
4. For all c \ne 0, a |b| if and only if ca |cb|.
```

```
Proof of (1)
     a | b 
\Rightarrow b = ak
\Rightarrow bc = ack
\Rightarrow bc = a(ck)
\Rightarrow a|bc
```
a "divides" b (a|b):

b = ak for some integer k

Simple Divisibility Facts

```
1. If a |b, then a |bc for all c.
2. If a \vert b and b \vert c, then a \vert c.
3. If a \vert b and a \vert c, then a \vert sb + tc for all s and t.
4. For all c \neq 0, a |b| if and only if ca |cb|.
```
Proof of (2)
\n
$$
a | b \Rightarrow b = ak_1
$$

\n $b | c \Rightarrow c = bk_2$
\n $\Rightarrow c = ak_1k_2$
\n $\Rightarrow a | c$

a "divides" b (a|b): **b = ak for some integer k**

Simple Divisibility Facts

```
1. If a |b, then a |bc for all c.
2. If a \vert b and b \vert c, then a \vert c.
3. If a \vert b and a \vert c, then a \vert sb + tc for all s and t.
4. For all c \neq 0, a |b| if and only if ca |cb|.
```

```
Proof of (3)
a \mid b \Rightarrow b = ak_1a \mid c \Rightarrow c = ak_2sb + tc= sak_1 + tak_2= a(sk_1 + tk_2)
```
a "divides" b (a|b):

b = ak for some integer k

 \Rightarrow a|(sb+tc)

This Lecture

- Direct proof
- Contrapositive
- Proof by contradiction
- Proof by cases

Proving an Implication

Goal: If P, then Q. (P implies Q)

Method 1: Write assume P, then show that Q logically follows.

Claim: \vert If r is irrational, then \sqrt{r} is irrational.

How to begin with?

What if I prove "If \sqrt{r} is rational, then r is rational", is it equivalent?

Yes, this is equivalent; proving "if P, then Q" is equivalent to proving "if not Q, then not P".

Rational Number

R is rational \Leftrightarrow there are integers a and b such that

Is 0.281 a rational number? Is 0 a rational number? Yes, 281/1000 Yes, 0/1

If m and n are non-zero integers, is (m+n)/mn a rational number?

Is the sum of two rational numbers a rational number? Yes, a/b+c/d=(ad+bc)/bd

Is x=0.12121212…… a rational number?

Note that 100x-x=12, and so x=12/99.

Proving an Implication

Goal: If P, then Q. (P implies Q)

Method 2: Prove the *contrapositive*, i.e. prove "not Q implies not P".

Claim: \vert If r is irrational, then \sqrt{r} is irrational.

Proof: \parallel We shall prove the contrapositive -"*if √r is rational, then r is rational*."

Since \sqrt{r} is rational, \sqrt{r} = a/b for some integers a,b.

So $r = a^2/b^2$. Since a,b are integers, a^2/b^2 are integers.

Therefore, r is rational. \Box Q.E.D.

 $(Q.E.D.)$ "which was to be demonstrated", or "quite easily done". \odot

Proving an "if and only if"

Goal: Prove that two statements P and Q are "**logically equivalent"**, that is, one holds if and only if the other holds.

Example:

An integer is even if and only if the its square is even.

Method 1: Prove P implies Q **and** Q implies P.

Method 1': Prove P implies Q **and** not P implies not Q.

Method 2: Construct a chain of if and only if statement.

Proof the Contrapositive

An integer is even if and only if its square is even.

Method 1: Prove P implies Q **and** Q implies P.

Statement: If m is even, then m^2 is even

Proof: m = 2k

$$
m^2 = 4k^2
$$

Statement: If m^2 is even, then m is even

Proof: $m^2 = 2k$ $m = J(2k)$

Proof the Contrapositive

An integer is even if and only if its square is even.

Method 1': Prove P implies Q **and** not P implies not Q.

Statement: If m^2 is even, then m is even Contrapositive: If m is odd, then m^2 is odd.

Proof (the contrapositive):

Since m is an odd number, m = 2k+1 for some integer k.

So m² =
$$
(2k+1)^2
$$

= $(2k)^2 + 2(2k) + 1$

So m^2 is an odd number.

This Lecture

- Direct proof
- Contrapositive
- Proof by contradiction
- Proof by cases

Proof by Contradiction

$\overline{P} \to \mathbf{F}$ $\frac{\rightarrow}{P}$

To prove P, you prove that not P would lead to ridiculous result, and so P must be true.

You are working as a clerk.

If you have won the lottery, then you would not work as a clerk.

You have not won the lottery.

Proof by Contradiction

Theorem: $\sqrt{2}$ is irrational.

Proof (by contradiction):

- Suppose $\sqrt{2}$ was rational.
- Choose *m*, *n* integers without common prime factors (always possible)

such that
$$
\sqrt{2} = \frac{m}{n}
$$

• Show that *m* and *n* are both even, thus having a common factor 2, a **contradiction**!

Proof by Contradiction

Theorem: $\sqrt{2}$ is irrational.

Proof (by contradiction): Want to prove both m and n are even.

Infinitude of the Primes

Theorem. There are infinitely many prime numbers.

Proof (by contradiction):

Assume there are only finitely many primes.

Let p_1 , p_2 , ..., p_N be all the primes.

We will construct a number N so that N is not divisible by any p_i .

By our assumption, it means that N is not divisible by any prime number.

On the other hand, we show that any number must be divisible by *some* prime.

It leads to a contradiction, and therefore the assumption must be false.

So there must be infinitely many primes.

Divisibility by a Prime

Theorem. Any integer n > 1 is divisible by a prime number.

•Let n be an integer.

- •If n is a prime number, then we are done.
- •Otherwise, $n = ab$, both are smaller than n .
- •If a or b is a prime number, then we are done.
- \cdot Otherwise, $a = cd$, both are smaller than a.
- •If c or d is a prime number, then we are done.
- •Otherwise, repeat this argument, since the numbers are getting smaller and smaller, this will eventually stop and we have found a prime factor of n.

Idea of induction.

Infinitude of the Primes

Theorem. There are infinitely many prime numbers.

Proof (by contradiction):

Let p_1 , p_2 , …, p_N be all the primes.

```
Consider p_1p_2...p_N + 1.
```
Claim: if p divides a, then p does not divide a+1.

Proof (by contradiction):

a = cp for some integer c a+1 = dp for some integer d \Rightarrow 1 = (d-c)p, contradiction because p>=2.

So none of $p_1, p_2, ..., p_N$ can divide $p_1p_2...p_N + 1$, a contradiction.

This Lecture

- Direct proof
- Contrapositive
- Proof by contradiction
- Proof by cases

Proof by Cases

 $p \vee q$ $p \rightarrow r$ $q \rightarrow r$ \ddot{r}

e.g. want to prove a nonzero number always has a positive square.

 x is positive or x is negative if x is positive, then $x^2 > 0$. if x is negative, then x^2 > 0. $x^2 \ge 0$.

The Square of an Odd Integer

$$
\forall \; \text{odd} \; n, \exists m,n^2=8m+1?
$$

Idea 0: find counterexample.

 $3^2 = 9 = 8 + 1$, $5^2 = 25 = 3 \times 8 + 1$ …… $131^2 = 17161 = 2145 \times 8 + 1$, ………

Idea 1: prove that $n^2 - 1$ is divisible by 8.

n ² – 1 = (n-1)(n+1) = ??…

Idea 2: consider (2k+1)²

 $(2k+1)^2$ = 4k²+4k+1

If k is even, then both k^2 and k are even, and so we are done.

If k is odd, then both k^2 and k are odd, and so k^2 +k even, also done.

Trial and Error Won't Work!

Fermat (1637): If an integer n is greater than 2,

then the equation $a^n + b^n = c^n$ has no solutions in non-zero integers a, b, and c.

Claim: $313(a^3 + b^3) = c^3$ has no solutions in non-zero integers a, b, and c.

False. But smallest counterexample has more than 1000 digits.

Euler conjecture:

 $h^4 + b^4 + c^4 = d^4$ has no solution for a,b,c,d positive integers.

Open for 218 years, until Noam Elkies found

 $95800^4 + 217519^4 + 414560^4 = 422481^4$

The Square Root of an Even Square

Statement: If m^2 is even, then m is even

Contrapositive: If m is odd, then m^2 is odd.

Proof (the contrapositive):

Since m is an odd number, m = 2l+1 for some natural number l.

So m² =
$$
(2I+1)^2
$$

= $(2I)^2 + 2(2I) + 1$

So m^2 is an odd number.

Proof by contrapositive.

Rational vs Irrational

Question: If a and b are irrational, can a^b be rational??

We know that $\sqrt{2}$ is irrational, what about $\sqrt{2^{2}}$?

Case 1: √2√2 is rational

Then we are done, a=√2, b=√2.

Case 2: √2√2 is irrational

Then $(\sqrt{2})^{\sqrt{2}} = \sqrt{2^2} = 2$, a rational number So a=**√2√2**, b= √2 will do.

So in either case there are a,b irrational and a^b be rational.

We don't (need to) know which case is true!

Summary

We have learnt different techniques to prove mathematical statements.

- Direct proof
- Contrapositive
- Proof by contradiction
- Proof by cases

Next time we will focus on a very important technique, proof by induction.